

From tidying books to shuffling cards

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Conférence MIPS
29 mai 2026

Outline

- 1 Tidying books
- 2 Shuffling cards
- 3 A touch of magic

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Tidying books



Tidying books

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Tidying books



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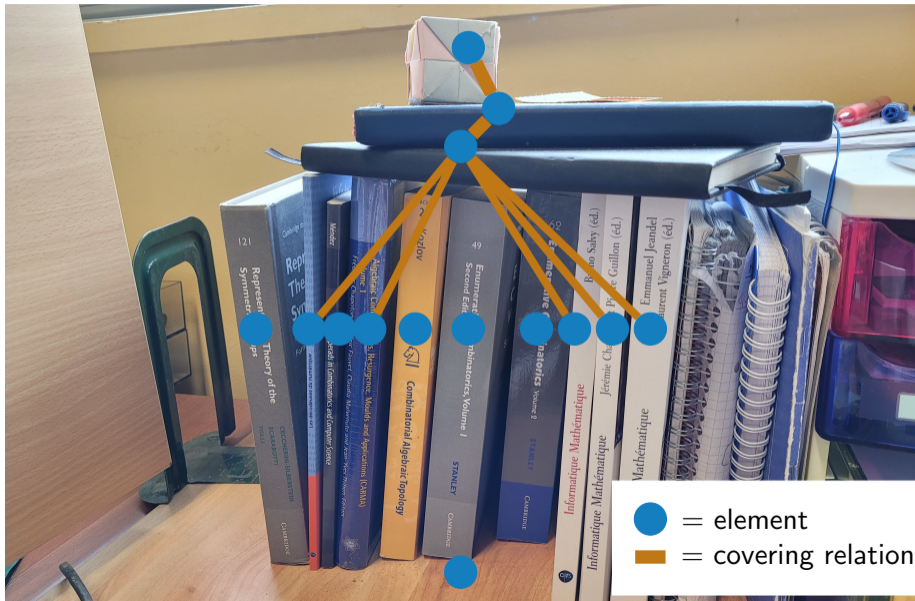


Tidying books



● = element
— = covering relations

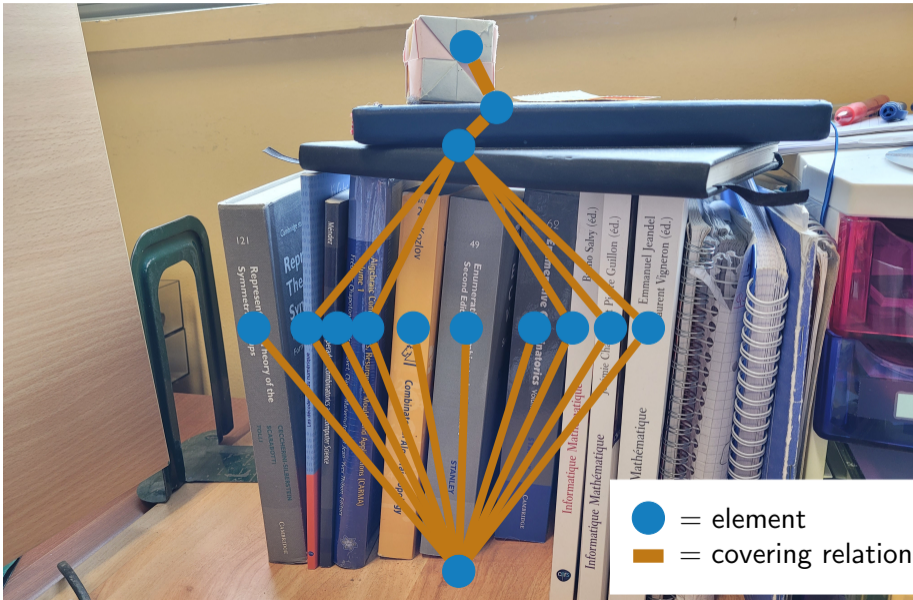
Tidying books



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Tidying books

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Sorting books

- On a shelf, the books are **totally ordered**



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- Here, we want to sort them in the alphabetic order



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Main questions :

- How many comparisons ? Is it efficient ?

Why do we need efficient sorting algorithms ?



Sorting books : The one-handed librarian's sort aka. Bubble sort 1

Algorithm

Going from left to right, the librarian

- compares two books
- exchanges them if needed

At the end of the shelf, starts again from the beginning

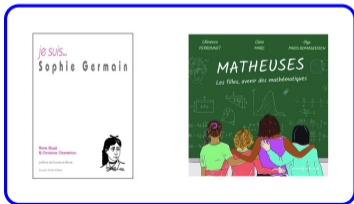
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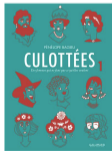
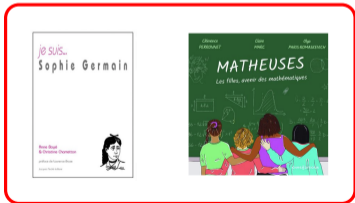
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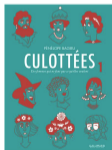
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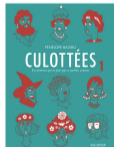
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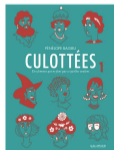
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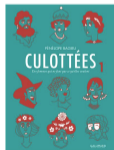
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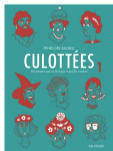
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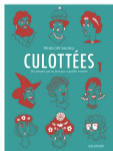
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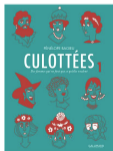
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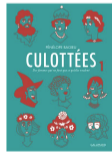
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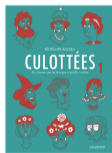
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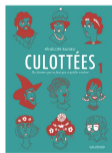
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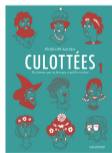
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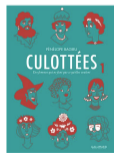
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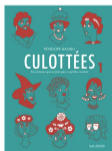
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→ Sorted !

Sorting permutations

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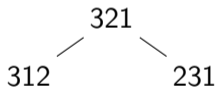
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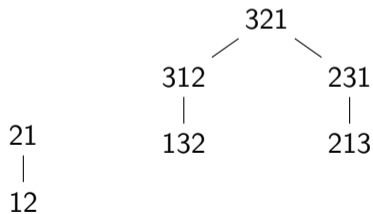


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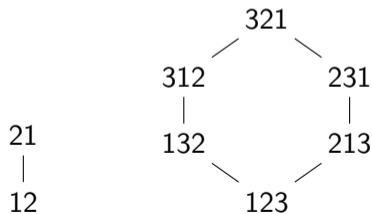
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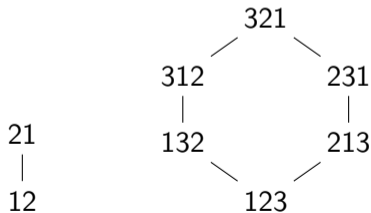
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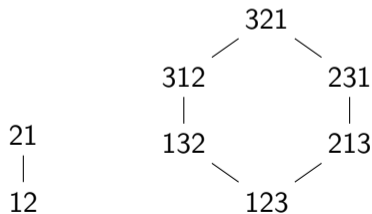
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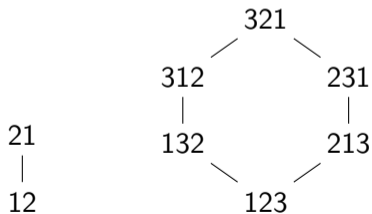
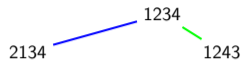
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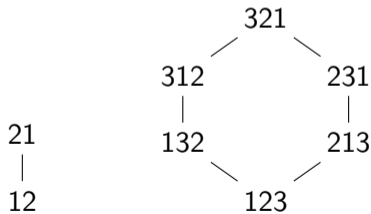
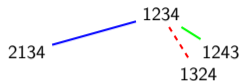


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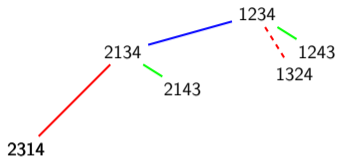
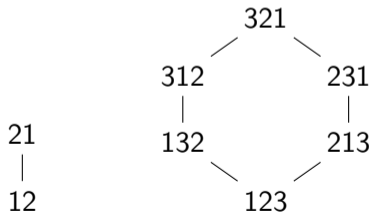


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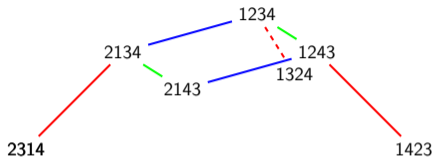
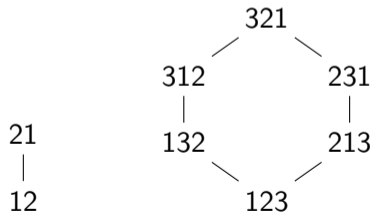
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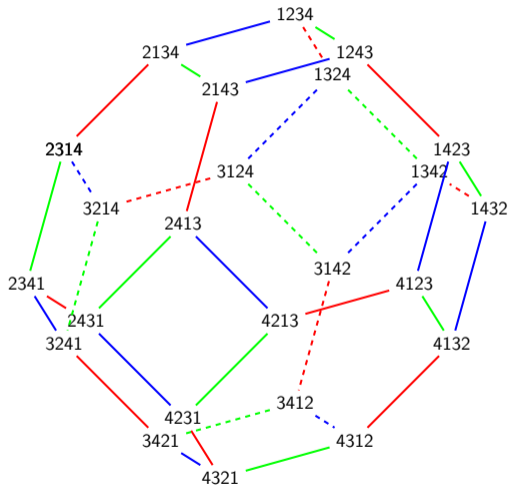
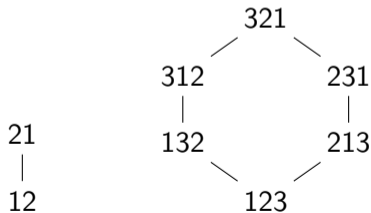


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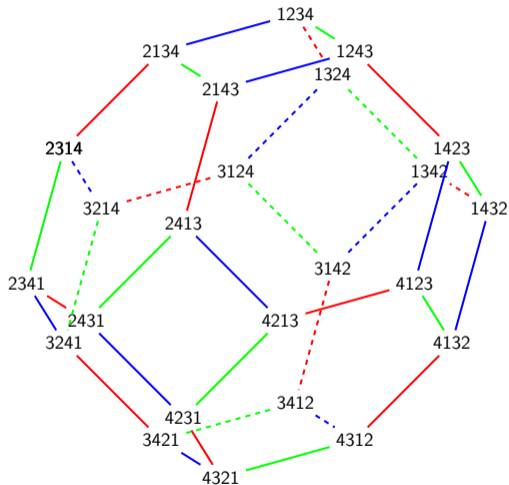
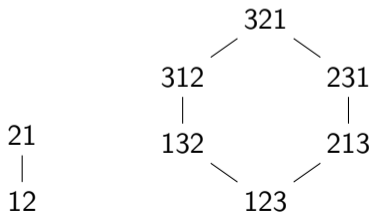
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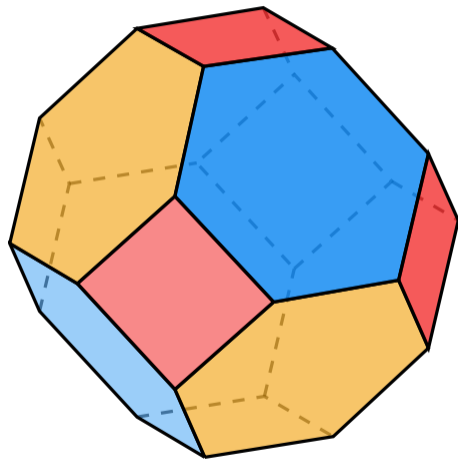
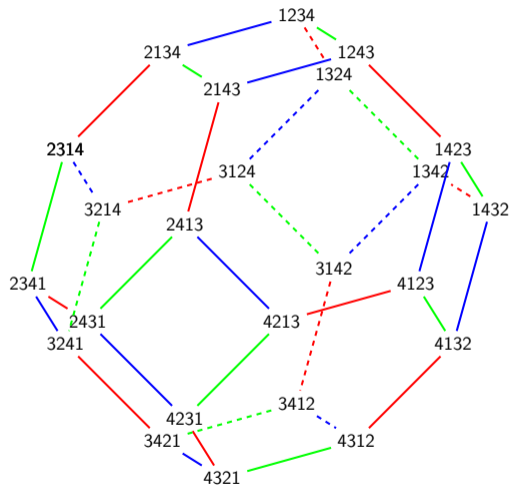
Sorting permutations

Answers [Friend, 1956]

- Does the order in which we do the exchange (from left to right) matter? No
- How many exchanges must be done? At most $\frac{n(n-1)}{2}$



Weak Bruhat order and permutohedron



Shuffling cards

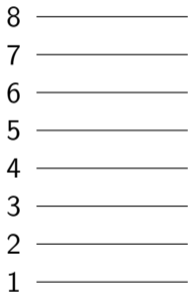
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- 2 Shuffling cards**
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Why shuffling cards ?

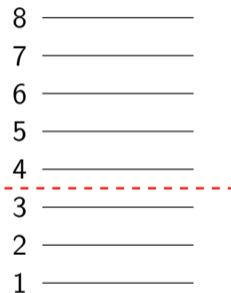


A naive (but useful) shuffle



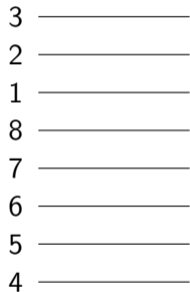
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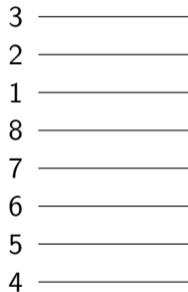
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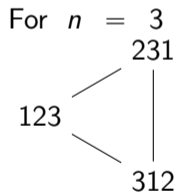


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How many
shuffles ?

n

How do you shuffle cards ?



More we shuffle cards, the better they are shuffled ?



<https://www.youtube.com/shorts/pwW6-YpWUuo>



Shuffles (in math)

Let \mathcal{A} be a finite alphabet and \mathcal{A}^* its set of (possibly empty) words (finite sequences of letters).

For any $a, b \in \mathcal{A}$ and $m, n \in \mathcal{A}^*$ the **shuffle product** is:

$$a.m \sqcup b.n = a.(m \sqcup b.n) + b.(a.m \sqcup n),$$

with $\varepsilon \sqcup m = m \sqcup \varepsilon = m$, where ε is the empty word.

Example :

$$ana \sqcup bob =$$

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Dovetail shuffle and probabilities

- Mixing the deck "1 2 3 4", we can get the following decks :

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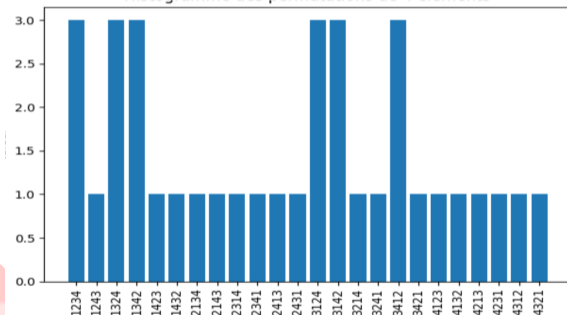
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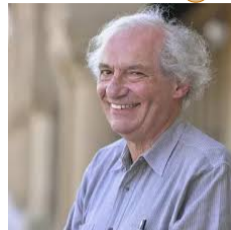


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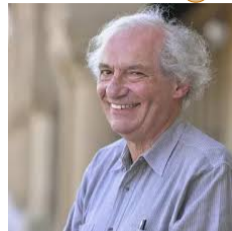
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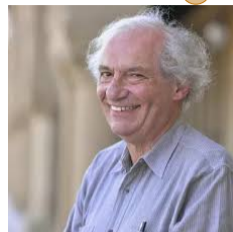


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How many shuffles to get a well-shuffled deck of 52 cards ?

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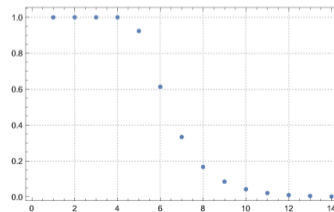


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TABLE 3
Total variation distance for m shuffles of 25, 32, 52, 78, 104, 208 or 312 distinct cards

m	1	2	3	4	5	6	7	8	9	10
25	1.000	1.000	0.999	0.775	0.437	0.231	0.114	0.056	0.028	0.014
32	1.000	1.000	1.000	0.929	0.597	0.322	0.164	0.084	0.042	0.021
52	1.000	1.000	1.000	1.000	0.924	0.614	0.334	0.167	0.085	0.043
78	1.000	1.000	1.000	1.000	1.000	0.893	0.571	0.307	0.153	0.078
104	1.000	1.000	1.000	1.000	1.000	0.988	0.772	0.454	0.237	0.119
208	1.000	1.000	1.000	1.000	1.000	1.000	1.000	0.914	0.603	0.329
312	1.000	1.000	1.000	1.000	1.000	1.000	1.000	0.999	0.883	0.565

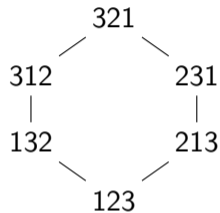


A touch of magic

Outline

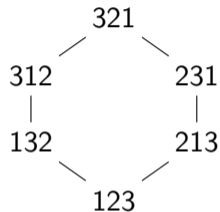
- 1 Tidying books
- 2 Shuffling cards
- 3 A touch of magic

Link between shuffling cards and tidying books



$$\left\{ \begin{array}{l} 1 \sqcup 23 = 123 + 213 + 231 \\ 1 \sqcup 32 = 132 + 312 + 321 \\ 12 \sqcup 3 = 123 + 132 + 312 \\ 21 \sqcup 3 = 213 + 231 + 321 \end{array} \right.$$

Link between shuffling cards and tidying books

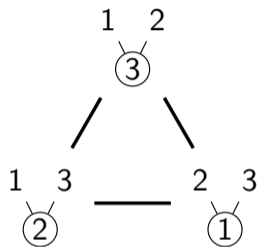


$$\begin{cases} 1 \sqcup 23 = 123 + 213 + 231 \\ 1 \sqcup 32 = 132 + 312 + 321 \\ 12 \sqcup 3 = 123 + 132 + 312 \\ 21 \sqcup 3 = 213 + 231 + 321 \end{cases}$$

Theorem [Loday-Ronco, 02]

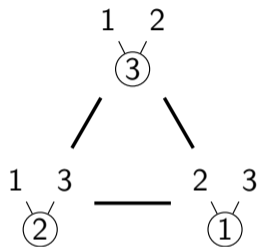
$$u \sqcup v = \sum_{(u \triangleleft v) \leq w \leq (u \triangleright v)} w$$

Could we shuffle something else ?



$$1 \sqcup (2 \sqcup 3) = 1 \sqcup \left(\begin{array}{c} 3 \\ \textcircled{2} \end{array} + \begin{array}{c} 2 \\ \textcircled{3} \end{array} \right) = 2 \times \begin{array}{c} 2 \ 3 \\ \textcircled{1} \end{array} + \begin{array}{c} 1 \ 3 \\ \textcircled{2} \end{array} + \begin{array}{c} 1 \ 2 \\ \textcircled{3} \end{array}$$

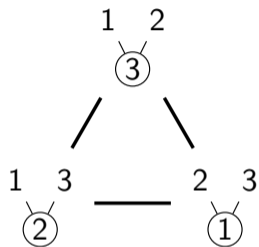
Could we shuffle something else ?



$$1 \sqcup (2 \sqcup 3) = 1 \sqcup \left(\begin{matrix} 3 \\ \textcircled{2} \end{matrix} + \begin{matrix} 2 \\ \textcircled{3} \end{matrix} \right) = 2 \times \begin{matrix} 2 & 3 \\ \textcircled{1} \end{matrix} + \begin{matrix} 1 & 3 \\ \textcircled{2} \end{matrix} + \begin{matrix} 1 & 2 \\ \textcircled{3} \end{matrix}$$

$$(1 \sqcup 2) \sqcup 3 = \left(\begin{matrix} 2 & 1 \\ \textcircled{1} & \textcircled{2} \end{matrix} \right) \sqcup 3 = \begin{matrix} 2 & 3 \\ \textcircled{1} \end{matrix} + \begin{matrix} 1 & 3 \\ \textcircled{2} \end{matrix} + 2 \times \begin{matrix} 1 & 2 \\ \textcircled{3} \end{matrix}$$

Could we shuffle something else ?

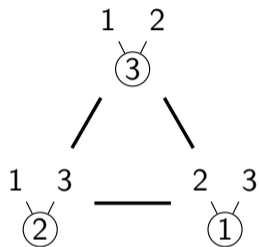


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\neq

$$(1 \sqcup 2) \sqcup 3 = \left(\begin{matrix} 2 & 1 \\ \textcircled{1} + \textcircled{2} \end{matrix} \right) \sqcup 3 = \begin{matrix} 2 & 3 \\ \textcircled{1} \end{matrix} + \begin{matrix} 1 & 3 \\ \textcircled{2} \end{matrix} + 2 \times \begin{matrix} 1 & 2 \\ \textcircled{3} \end{matrix}$$

Could we shuffle something else ?



$$1 \sqcup (2 \sqcup 3) = 1 \sqcup \left(\begin{matrix} 3 \\ \textcircled{2} \end{matrix} + \begin{matrix} 2 \\ \textcircled{3} \end{matrix} \right) = 2 \times \begin{matrix} 2 & 3 \\ \textcircled{1} \end{matrix} + \begin{matrix} 1 & 3 \\ \textcircled{2} \end{matrix} + \begin{matrix} 1 & 2 \\ \textcircled{3} \end{matrix}$$

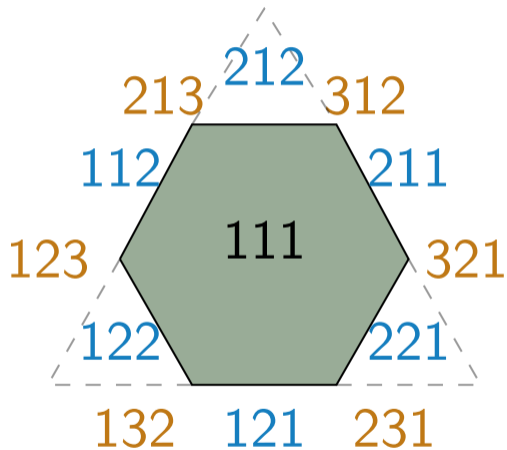
\neq

$$(1 \sqcup 2) \sqcup 3 = \left(\begin{matrix} 2 & 1 \\ \textcircled{1} & \textcircled{2} \end{matrix} \right) \sqcup 3 = \begin{matrix} 2 & 3 \\ \textcircled{1} \end{matrix} + \begin{matrix} 1 & 3 \\ \textcircled{2} \end{matrix} + 2 \times \begin{matrix} 1 & 2 \\ \textcircled{3} \end{matrix}$$

Idea :

- Shuffle faces (and not just vertices)
- Of "good" nestohedra

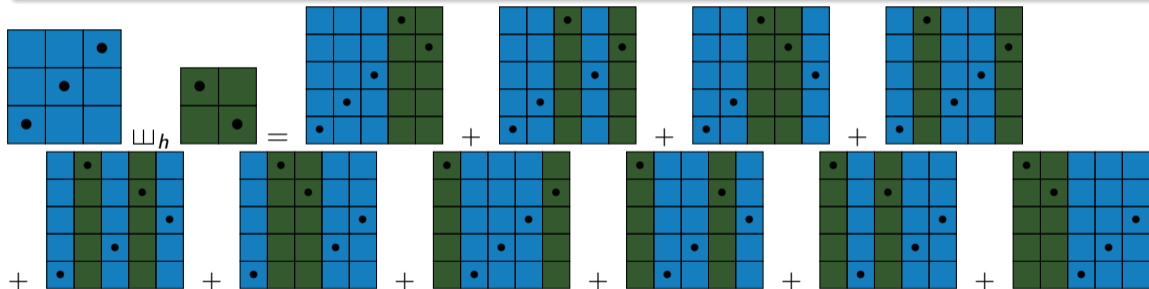
Let's shuffle faces !



Back to the previous shuffle on permutations : Horizontal shuffle ● ● ● 3

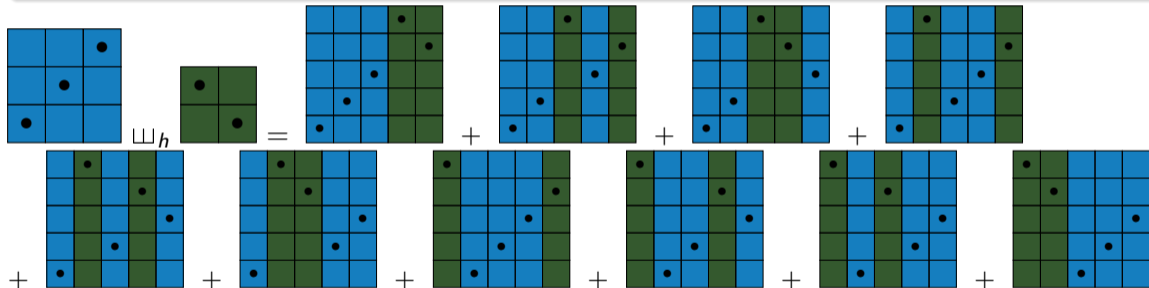
Example:

$$123 \sqcup_h 21 = 12354 + 12534 + 12543 + 15234 + 15243 + 15423 + 51234 + 51243 + 51423 + 54123$$



Example:

$$123 \sqcup_h 21 = 12354 + 12534 + 12543 + 15234 + 15243 + 15423 + 51234 + 51243 + 51423 + 54123$$



Question

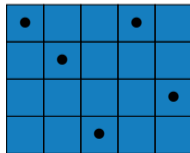
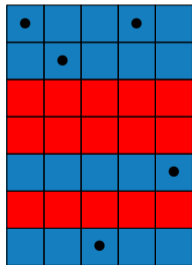
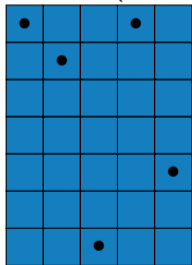
How could we get something like 12321 ?

Vertical shuffle

For σ and τ two surjections,

$$\sigma \sqcup_v \tau = \sum_{\substack{\text{std}(s)=\sigma \\ \text{std}(t)=\tau}} s.t,$$

where $\text{std}(76173)=43132$



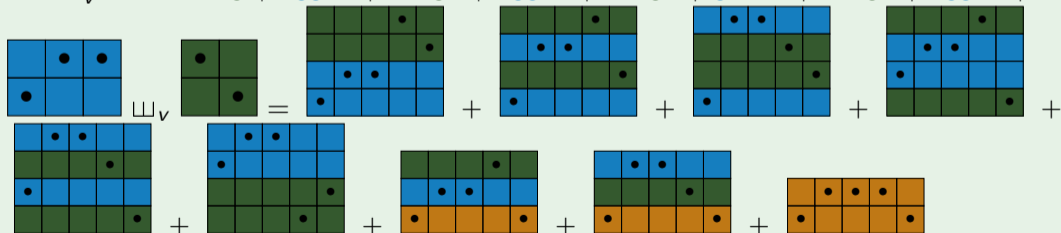
Vertical shuffle

For σ and τ two surjections,

$$\sigma \sqcup_v \tau = \sum_{\substack{\text{std}(s)=\sigma \\ \text{std}(t)=\tau}} s.t,$$

Example:

$$122 \sqcup_v 21 = 12243 + 13342 + 14432 + 23341 + 24431 + 34421 + 12231 + 13321 + 12221$$



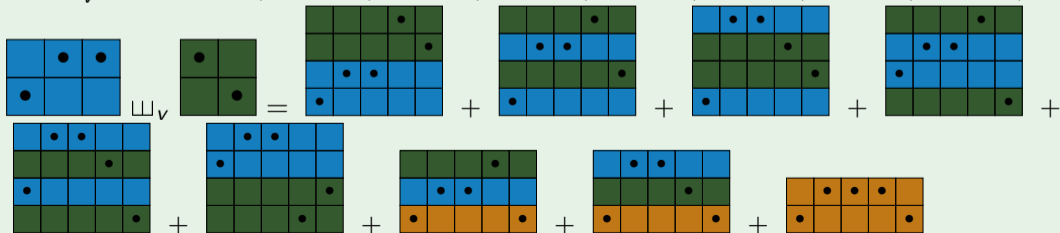
Vertical shuffle: tridendriform products

For σ and τ two surjections,

$$\sigma <_v \tau = \sum_{\substack{\text{std}(s)=\sigma \\ \text{std}(t)=\tau \\ \min(s,t)=\min(s)}} s.t, \quad \sigma >_v \tau = \sum_{\substack{\text{std}(s)=\sigma \\ \text{std}(t)=\tau \\ \min(s,t)=\min(t)}} s.t, \quad \sigma \cdot_v \tau = \sum_{\substack{\text{std}(s)=\sigma \\ \text{std}(t)=\tau \\ \min(s,t)=\min(s)=\min(t)}} s.t,$$

Example:

$$122 \sqcup_v 21 = 12243 + 13342 + 14432 + 23341 + 24431 + 34421 + 12231 + 13321 + 12221$$



Tridendriform algebras

Definition (Loday, Ronco, 2004 ; Chapoton 2002)

A tridendriform algebra is a vector space A endowed with products $\langle : A \otimes A \rightarrow A$, $\cdot : A \otimes A \rightarrow A$ and $\rangle : A \otimes A \rightarrow A$, such that:

- 1 $(a \langle b) \langle c = a \langle (b * c)$,
- 2 $(a * b) \rangle c = a \rangle (b \rangle c)$,
- 3 $(a \rangle b) \langle c = a \rangle (b \langle c)$,
- 4 $(a \cdot b) \cdot c = a \cdot (b \cdot c)$,
- 5 $(a \rangle b) \cdot c = a \rangle (b \cdot c)$,
- 6 $(a \langle b) \cdot c = a \cdot (b \rangle c)$,
- 7 $(a \cdot b) \langle c = a \cdot (b \langle c)$,

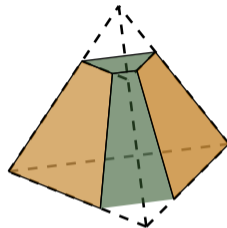
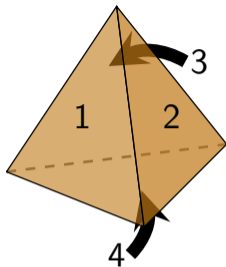
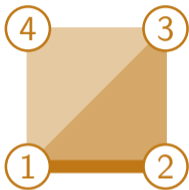
with $* = \langle + \cdot + \rangle$

Theorem

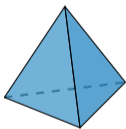
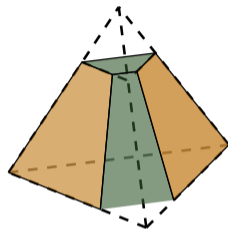
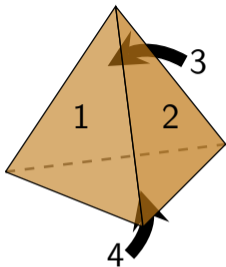
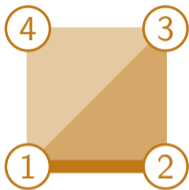
[Burgunder-Curien-Ronco, 15]

The tridendriform algebra of surjections is free.

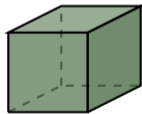
Nestohedra



Nestohedra



Simplices



Hypercubes



Associahedra



Permutohedra

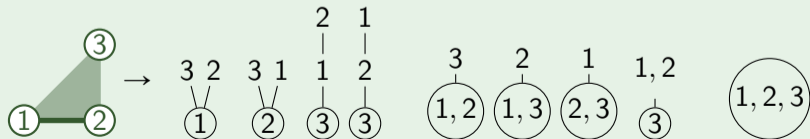
Constructs

A **construct** c of a hypergraph H is a rooted tree whose vertices form a partition of $V(H)$ and defined inductively by:

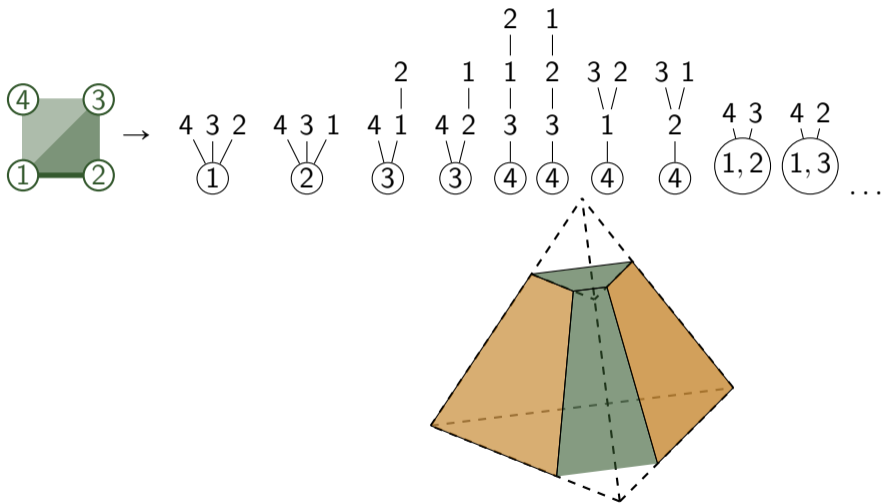
- c has only one node labelled by $V(H)$,
- or the root of c is $E \subseteq V(H)$
and each of its children is a construct of a connected component of $H - E$.

The set of constructs of a given hypergraph labels faces of the associated polytope.

First example:



Second example geometrically



Heuristics for a tridendriform structure

Let $\mathbf{H}^{\mathcal{X}}$ be a family of hypergraph polytopes, indexed by some finite sets \mathcal{X} (sets of vertices of the associated hypergraphs).

For $S = A(S_1, \dots, S_m)$ and $T = B(T_1, \dots, T_n)$ two constructs of $\mathbf{H}^{\mathcal{X}}$ and $\mathbf{H}^{\mathcal{Y}}$ respectively (\mathcal{X}, \mathcal{Y} disjoint), we would like to define the following operations

- $S < T$ as a sum of constructs of $\mathbf{H}^{\mathcal{X} \cup \mathcal{Y}}$ having **root** A ,
- $S > T$ as a sum of constructs of $\mathbf{H}^{\mathcal{X} \cup \mathcal{Y}}$ having **root** B ,
- $S \cdot T$ as a sum of constructs of $\mathbf{H}^{\mathcal{X} \cup \mathcal{Y}}$ having **root** $A \cup B$.

Heuristics for a tridendriform structure

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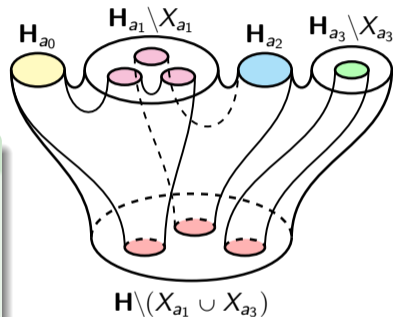
First problem in hypercubes

$$1 * (2 * 3) = 1 * \left(\begin{array}{c} 3 \\ | \\ \textcircled{2} \end{array} + \begin{array}{c} 2 \\ | \\ \textcircled{3} \end{array} + \textcircled{23} \right) = 3 \times \begin{array}{c} 3 \ 2 \\ \vee \\ \textcircled{1} \end{array} + \dots \neq (1 * 2) * 3$$

A team is a (resp. quasi-strict) strict team if each edge of a hypergraph \mathbf{H}_a is connected in \mathbf{H} (resp. or totally disconnected).

Examples:

- Strict teams : Associahedra, Permutohedra, Restrictohedra, ...
- Quasi-strict teams : Simplices, Hypercubes, Erosohedra, ...



$$(X_{a_0} = X_{a_2} = \emptyset)$$

Shuffle product

Considering a team E and denoting by δ a tuple of constructs of the team's participating hypergraphs, we inductively associate to δ a sum of constructs of the supporting hypergraph:

$$*(\delta) = \sum_{\emptyset \subset B \subseteq A} q^{|B|-1} *_{B}(\delta),$$

where

$$*_{B}(\delta) = \left(\bigcup_{b \in B} X_b \right) (*(\delta_1^B), \dots, *(\delta_{n_B}^B)).$$

Proposition (Curien-D.O.-Obradovic, 25)

$$*_{B}(\delta) = \sum_{U: \mathbf{H}, \text{root}(U)=X_B \text{ and } \forall a \in A, U|_{\mathbf{H}_a} = C_a} q^{\mu^{\tau}(U) - |B| + 1} U.$$

Theorem [**BDO**–Curien–Obradović, 25 and 26+]

- q -tridendriform product on faces of "strict" nestohedra (including associahedra, permutohedra, restrictohedra, friesohedra...) + link with an order
- (-1) -tridendriform product on faces of "semi-strict" nestohedra (including simplices, hypercubes, erosohedra...)

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Back to the simplex:

$$\begin{aligned}
 1 \sqcup (2 \sqcup 3) &= 1 \sqcup \left(\begin{array}{c} 3 \\ \textcircled{2} \end{array} + \begin{array}{c} 2 \\ \textcircled{3} \end{array} - \textcircled{2 \ 3} \right) \\
 &= \begin{array}{c} 2 \ 3 \\ \textcircled{1} \end{array} + \begin{array}{c} 1 \ 3 \\ \textcircled{2} \end{array} + \begin{array}{c} 1 \ 2 \\ \textcircled{3} \end{array} - \textcircled{1 \ 2} - \textcircled{1 \ 3} - \textcircled{2 \ 3} + \textcircled{1 \ 2 \ 3}
 \end{aligned}$$

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 \end{aligned}$$

Thank you !